Secure Applications without Secure Infrastructures

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Critical infrastructures

- We have to come to rely on the Internet to an extent that life without it becomes difficult to image.
  - Air travel: no more paper tickets, only e-tickets since 2008; booking via web sites.
  - Conference registration: via web sites
  - Payment: credit card details entered on web sites; PayPal.
  - Communication: via email
  - Plus e-banking, e-commerce, e-government, SCADA, …
- Internet & web have become critical infrastructures.

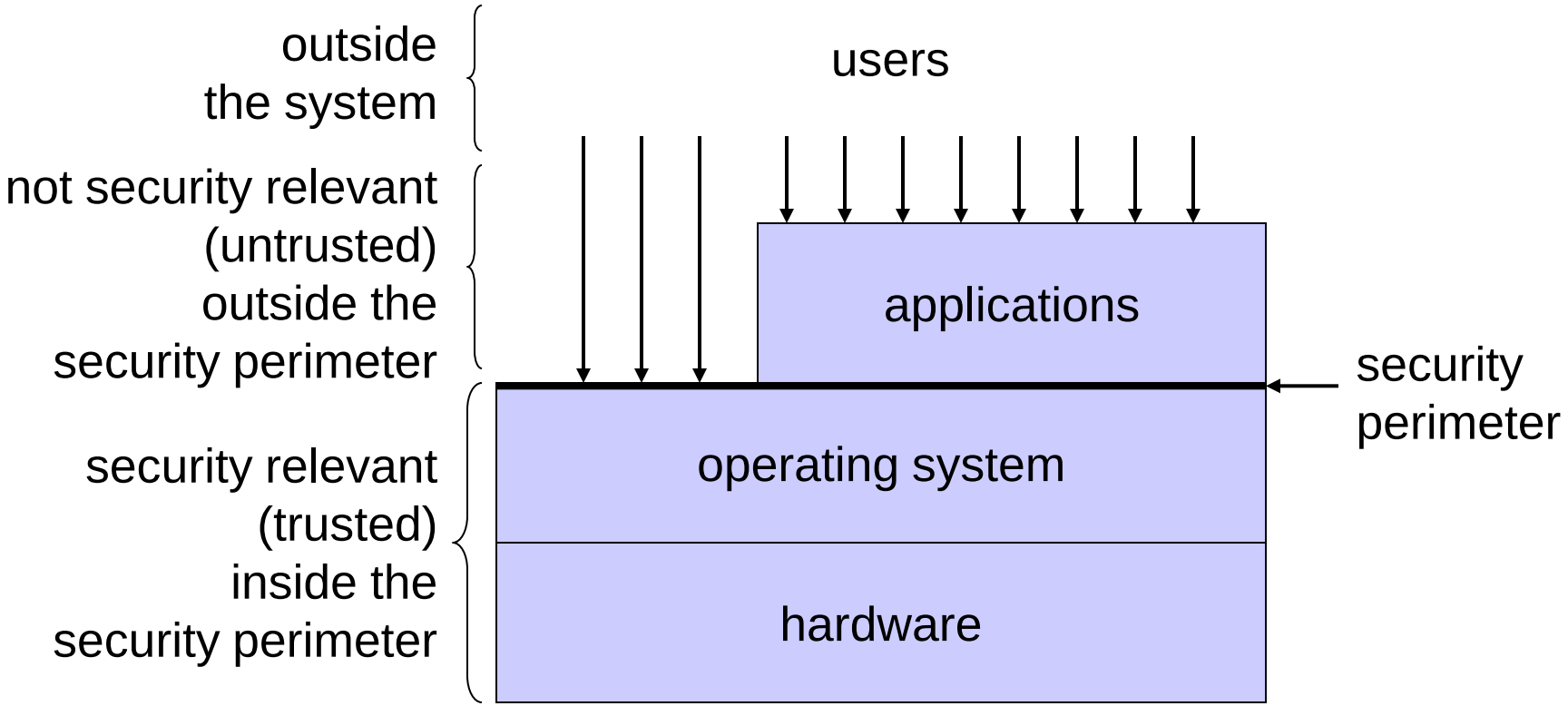
Do we have to secure this critical infrastructure?

Infrastructure security

- Historically, computer and communications security are infrastructure security.
- Computer security = operating system security:
  O/S is the infrastructure for users and applications.
  - Provides process isolation, access control, …
  - Once data are with the application the job is done.
- Communications security = secure channels:
  infrastructure carrying data from sender to receiver.
  - Once data are with the receiver the job is done.

Computer security, 1988

Morrie Gasser, Building a Secure Computer System, Van Nostrand Reinhold



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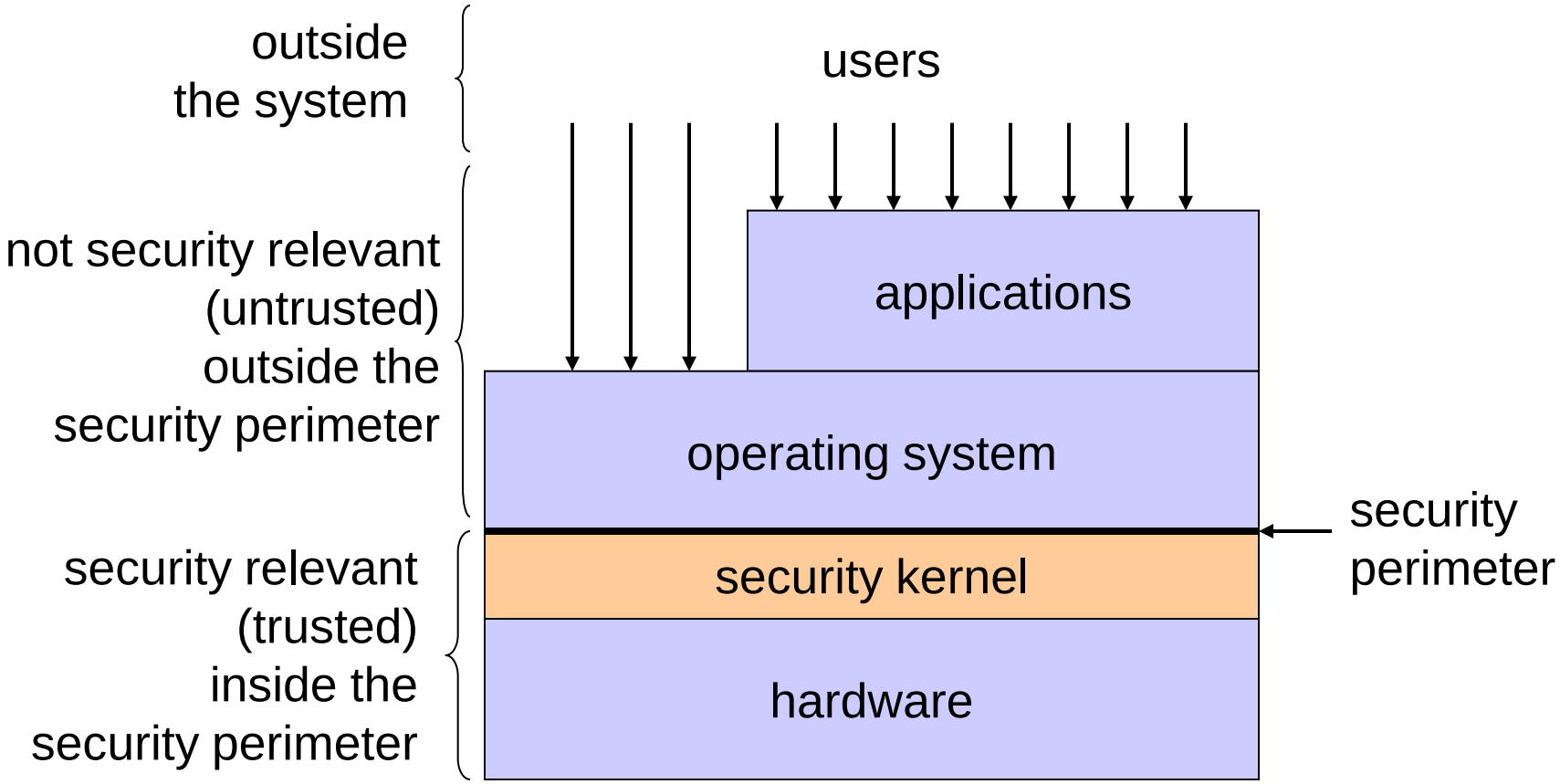


Security strategy – the past

- Higher security by small, verifiable security kernel that implements the reference monitor.
- Security guaranteed by the lower system layers (the infrastructure).
- Applications need not be trusted.
- Security evaluation (Orange Book): focus on infrastructures managed by IT professionals.
- The defenders retreat to the security kernel.

Communications security

- Focus on the design of secure channels: IPsec, IPsec over IPsec, SSL/TLS, …
- Infrastructure services at network and transport layer.
- Protect against attackers (“spies”) who can read, modify, delete, insert, replay messages.
- Job done once messages are delivered.
- No protection against attacks in the end systems (“hackers”).

Secure channels

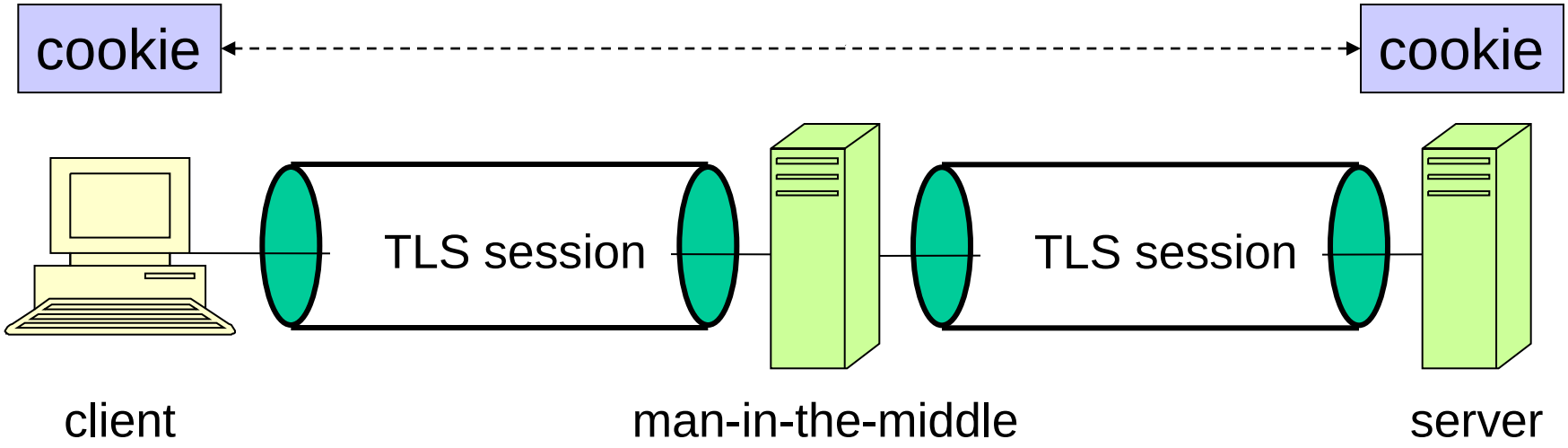
- Formal security proofs for some protocols exist (TLS).
- Challenge: combining channels at different layers.
- Attacks become possible when there is a mismatch of channel end points.
- Tunnels at different layers by definition have different end points!
- ‘Server’ and ‘client’ are dangerous simplifications.

Case study: TLS & e-commerce

- Server is authenticated to client when setting up the TLS tunnel (unilateral authentication).
- User has account at server, protected by password.
- Password based authentication within SSL tunnel.
  - Attacker cannot sniff user password.
- Server authenticates user, returns HTTP cookie.
- Client browser includes cookie in further requests to server; requests are attributed to authenticated user.

Man-in-the-middle attack

Attacker splices two TLS sessions.



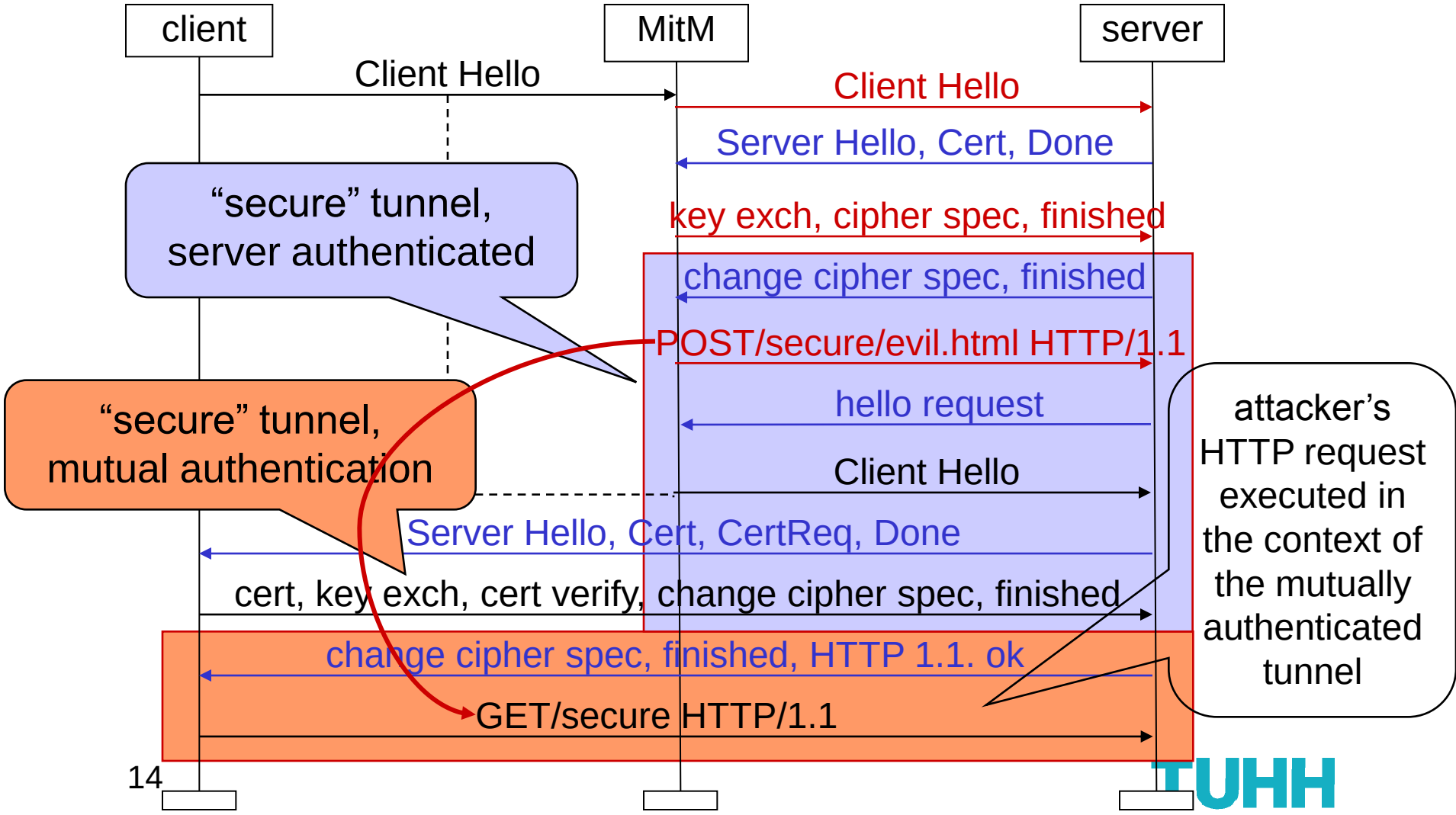
As a defence, bind cookie to TLS session.

Recent TLS security scare

- “Flaw” of TLS widely reported.
  - Marsh Ray, Steve Dispensa: Renegotiating TLS, 4.11.2009
- Background: TLS employed for user authentication when accessing a secure web site.
- Common practice for web servers to let users start with an anonymous TLS session.
- Request for a protected resource triggers TLS renegotiation; mutual authentication requested when establishing the new TLS tunnel.

Bugtraq ID 36935: “Multiple vendors’ TLS protocol implementations are prone to a security vulnerability related to the session-renegotiation process.”

Recent https-Problem



Comment

- Application developers using session renegotiation for user authentication made assumptions about renegotiation I failed to spot in RFC 5246.
- Fact: typical use case for renegotiation suggests that the new session is a continuation of the old session.
  - Plausible assumptions about a plausible use case are treated as a specification of the service.
- Fix: TLS renegotiation cryptographically tied to the TLS connection it is performed in (RFC 5746).
  - TLS adapted to meet expectations of application.
- This had really been an application layer problem.
  - State at server persists over two TLS tunnels; attacker sends a malicious partially complete command in the first tunnel.

CompSec & CommSec – today

- Cross-Site Scripting
  - #1 in 2007 OWASP Top Ten Vulnerabilities
  - #1 in CVE seit 2005
- SQL Injection
  - #2 in CVE seit 2006
- Cross-Site Request Forgery
- JavaScript Hijacking
- DNS Rebinding
- Code injection attacks,
- …

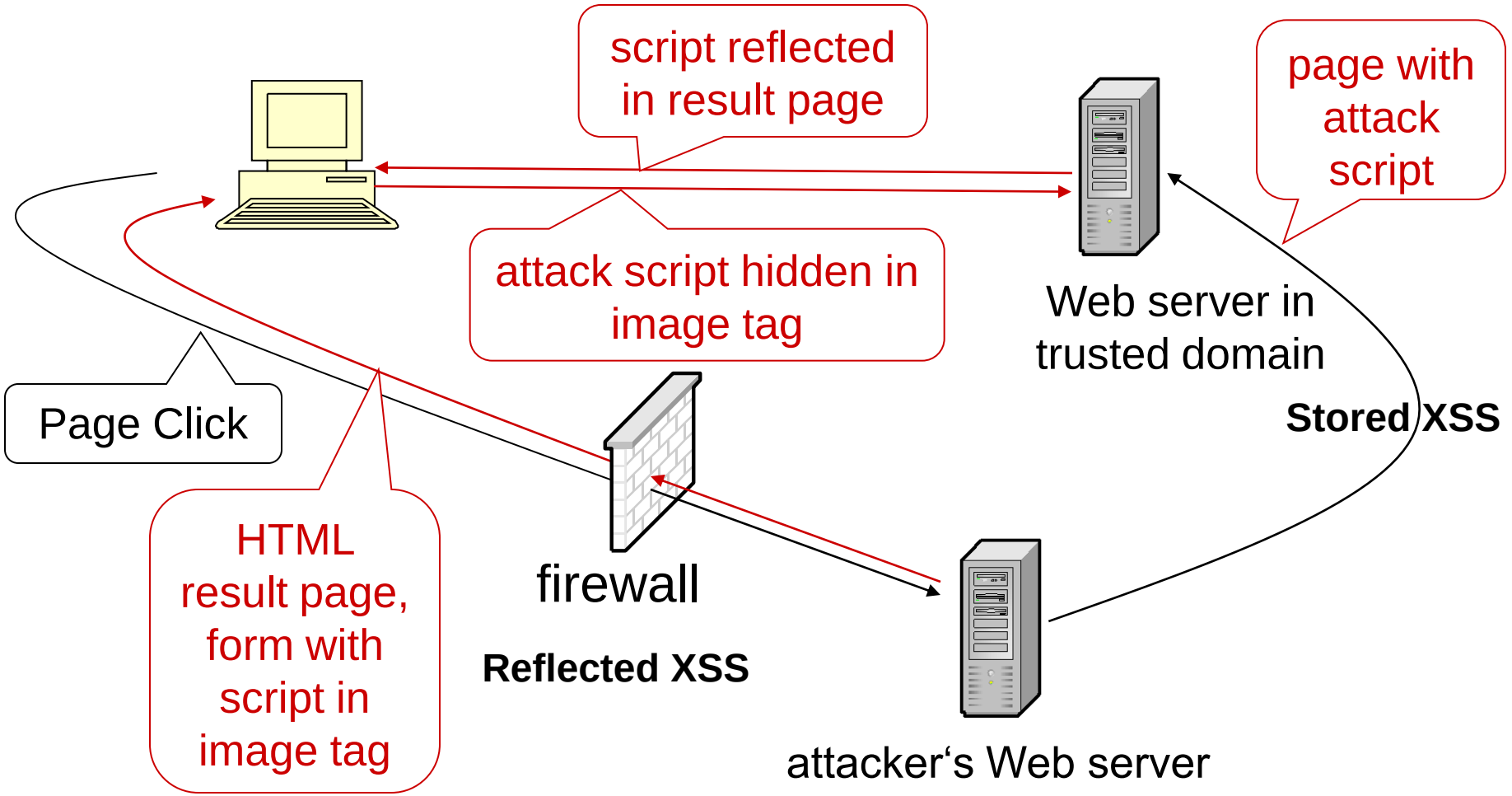
Client-side scripting (Web 1.0)

- Response page may contain scripts (often written in JavaScript) that will be executed in browser.
- Attack might be launched by placing malicious script in a response page.
- Browsers enforce a JavaScript same origin policy:
  - Script may get access to its own DOM only.
  - Script may only connect to the DNS domain it came from.

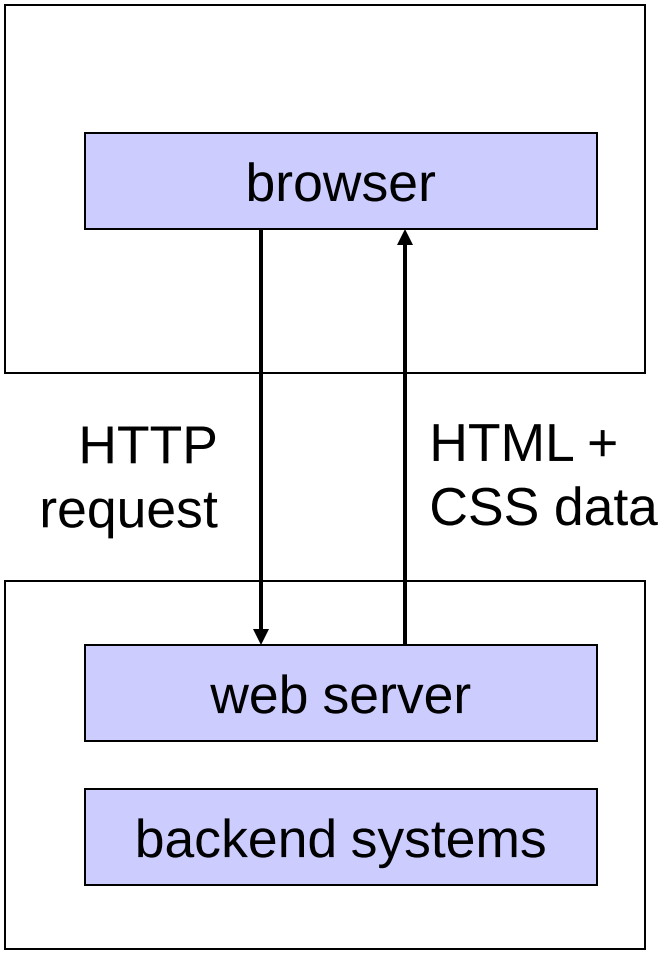
Cross Site Scripting – XSS

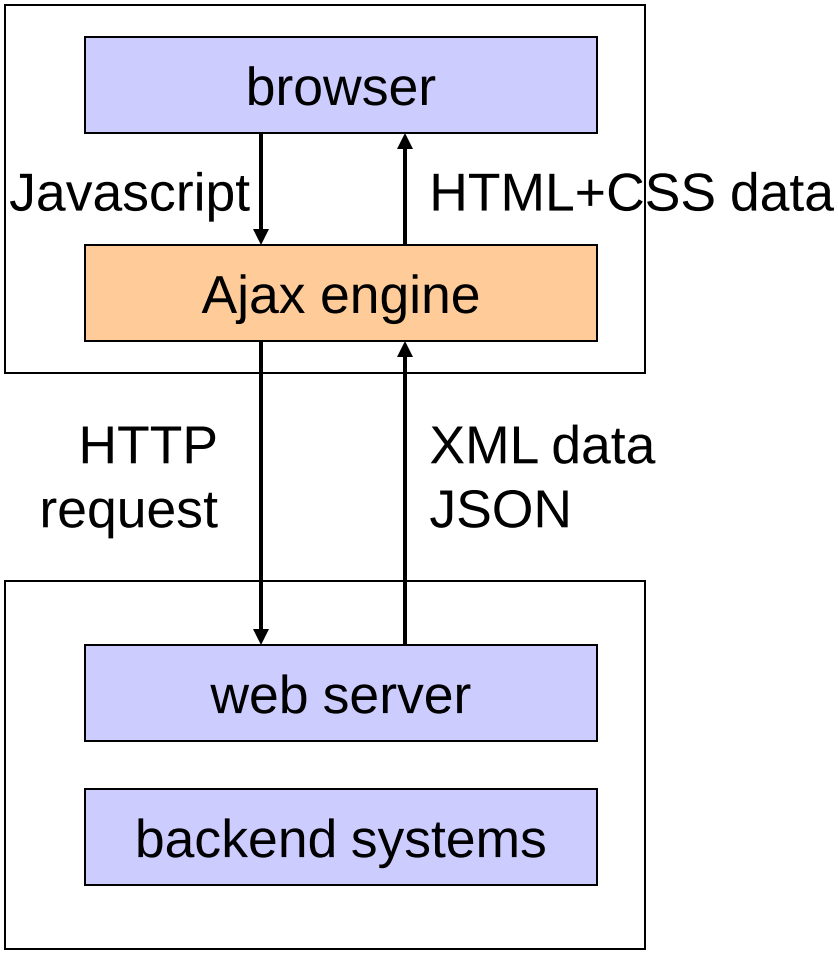
- Elevation of privilege attack involving attacker, client (victim), server ('trusted' by client).
  - Trust: code in pages from server executed with higher privileges at client (origin based access control).
- Attacker places malware on a page at server (stored XSS) or in a hidden form in a page on the attacker's host (reflected XSS).
  - Stored XSS: e.g. via a bulletin board.
  - Reflected XSS: e.g. via a search term.
- Malware returned by server to client in result page; executed at client with permissions of trusted server.

Cross-site scripting



Web 1.0 & Web 2.0


browser
HTTP
request
HTML +
CSS data
web server
backend systems


browser
Javascript
HTML+CSS data
Ajax engine
HTTP
request
XML data
JSON
web server
backend systems

JavaScript hijacking (Web 2.0)

- Exploits that a client side Ajax engine sitting between browser and web server performs many actions automatically.
- Exploits the fact that Web 2.0 applications may use JavaScript (JSON) for data transport.
- JSON string is a serialized JavaScript object, turned back into an object with JavaScript by calling eval() with the JSON string as the argument using the JavaScript object constructor.
- Attack rewrites constructor to disclose confidential data to attacker; bypasses same origin policy.

DNS rebinding

- Same origin policy: Script can only connect back to the server it was downloaded from.
- To make a connection, the client's browser needs the IP address of the server.
- Authoritative DNS server resolves DNS names in its domain to IP addresses.
- The client's browser 'trusts' the DNS server when enforcing the same origin policy.
- Trust is Bad for Security!

DNS rebinding attack

- Client visits attacker.org; attacker's DNS server resolves this name to attacker's IP address with short time-to-live.
- Attack script waits before connecting to attacker.org.
- Binding at browser has expired; new request for IP address of attacker.org, now bound to target address.
- Defence: Don't trust the DNS server on time-to-live; pin host name to original IP address;
  - J. Roskind: Attacks against the Netscape browser. in RSA Conference, April 2001.

DNS rebinding attack

- More sophisticated authorisation system: client browser refers to policy obtained from DNS server when deciding on connection requests.
- A malicious DNS server can thus authorize connection to the victim.
- Defence: Double check policy with the host at the IP address the DNS name is being resolved to.
  - Related to reverse DNS lookup.
  - Similar defences against bombing attacks in network security.
- Digital signatures do not help against DNS rebinding.

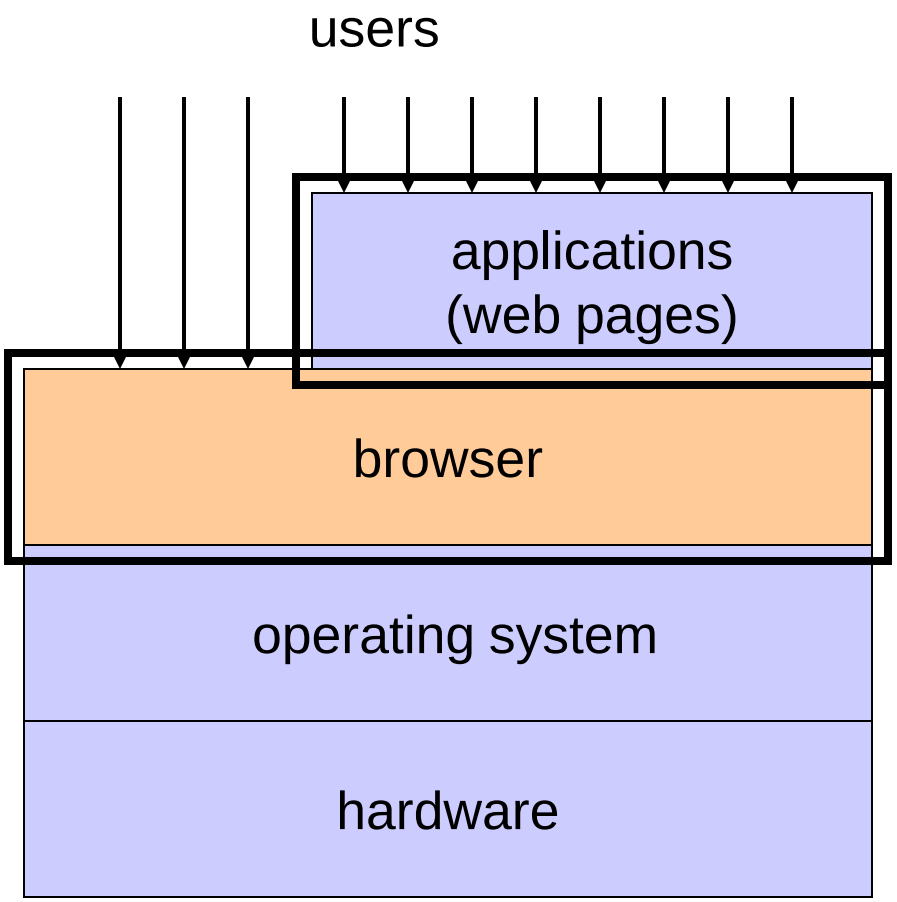
Computer security – today

- Same Origin Policies: who may read cookies, where may a script connect to?
- Reference monitor in the browser (sandbox).
- The attacks just sketched exploit vulnerabilities in this reference monitor.

- Web pages filter inputs, e.g. to defend against SQL injection attacks.
- Reference monitor moves to the application layer; application developers have to include relevant filters.

Computer security – today

security perimeter?



Summary

- Mechanisms in the traditional security kernel hardly defend against today's new attacks.
- Traditional secure channels hardly defend against today's new attacks.
- The line of defence against current attacks moves up to the application layer.
- Security mechanisms are moving out of the infrastructure into the applications.
- Defenders meet the attackers in front of the gates.

Research challenges

- Access control – mechanisms: modelling the mechanisms in the browser
  - considering Web 2.0, plug-ins, mashups, …
  - considering new mechanisms for authenticating data origin.
- Access control – policies: specification and enforcement of Cross Domain Policies.
  - Ajax Cross Domain Policies
  - HTTP Access Control Headers for Cross-Domain Policies, http://www.w3.org/TR/access-control/
  - Re-evaluating the role of the end user in setting policies.

Research challenges

- Authentication: examining necessary conditions on ‘authentication’.
  - Authentication can go beyond ‘corroborating the identity of the sender.
  - Would ‘recognition’ or ‘know thyself’ suffice?
- Interaction between protocol layers: understanding how to build tunnels in tunnels.
  - Re-evaluating which security services should be provided by the lower layers (the infrastructure) and which are provided within the application.

- We do not have to secure the infrastructure but the critical applications.
- Securing the critical infrastructure is neither sufficient nor necessary.
- Thank you very much for your attention.